Multi-party sessions as a security protocol abstraction

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Secure distributed programming

Only realistic security assumption:

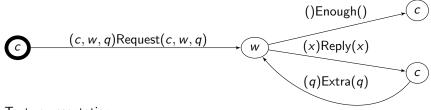
The network and any coalition of peers are potentially malicious.

Designing a (correct) security protocol by hand is hard:

- involves low-level, error-prone coding below communication abstractions,
- depends on global message choreography,
- needs to protect against coalitions of compromised peers.

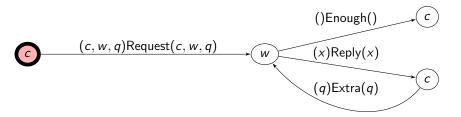
Therefore, our solution:

- to automatically generate tailored cryptographic protocols protecting against the network and compromised peers;
- to hide implementation details and provide mechanised proofs of correctness.

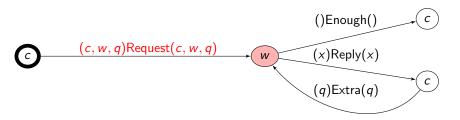


Text representation:

```
protocol WSn(role c, role w) {
    {c,w,q} Request {c,w,q} from c to w;
    rec loop {
      choice at w {
        {x} Reply {x} from w to c;
        {q} Extra {q} from c to w;
      continue loop; }
    or { Enough from w to c; } }
```



Execution Labels: Store: c: w: q: x:



Execution

Labels: Request

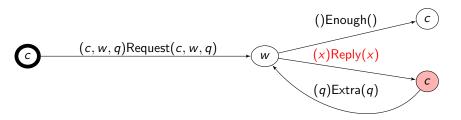
Store:

c: Alice

w: Bob

"Gone with the wind"

q: *x*:



Execution

Labels: Request-Reply

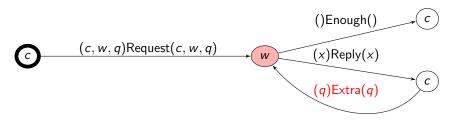
Store:

c: Alice

w: Bob

q: "Gone with the wind"

x: "8 euros"



Execution

Labels: Request-Reply-Extra

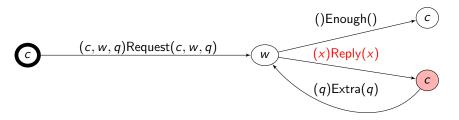
Store:

c: Alice

w: Bob

q: "In stock?"

x: "8 euros"



Execution

Labels: Request-Reply-Extra-Reply

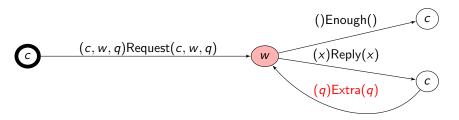
Store:

c: Alice

w: Bob

q: "In stock?"

x: "yes"



Execution

Labels: Request-Reply-Extra-Reply-Extra

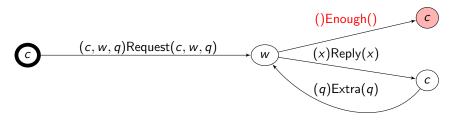
Store:

c: Alice

w: Bob

q: "Delivery date?"

x: "yes"



Execution

Labels: Request-Reply-Extra-Reply-Extra-Enough

Store:

c: Alice

w: Bob

q: "Delivery date?"

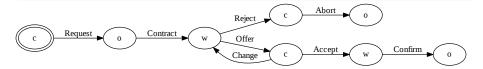
x: "yes"

Threats against session integrity

Powerful Attacker model

- can spy on transmitted messages
- can join a session as any role
- can initiate sessions

- can access the libraries (networking, crypto)
- cannot forge signatures



Attacks against an insecure implementation

- (Integrity) Rewrite Offer by Reject
- (Replay) Intercept Reject and replay old Offer, triggering a new iteration
- (Sender authentication) Intercept Abort and send Confirm to o
- ... and many more against the store

Protocol outline

Principles of our protocol generation

- Each edge is implemented by a unique concrete message.
- We want static message handling for efficiency.

Against replay attacks

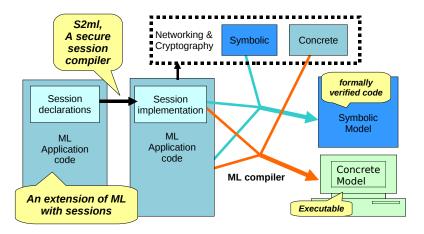
- between session executions: session nonces
- between loop iterations: time stamps
- at session initialisations: anti-replay caches



Against session flow attacks

• Signatures of the entire message history (optimisations possible ...)

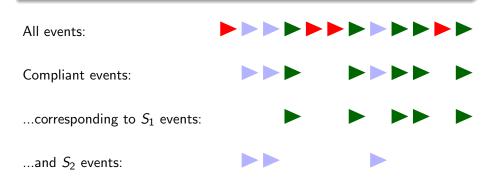
Architecture



Security result

Theorem (Session Integrity)

For any run of a $S_1....S_n$ -system, there is a partition of the compliant events such that each equivalence class coincides with a compliant subtrace of a session S_i from from $S_1...S_n$.



Conclusion

- Security protocols are hard to write by hand. They are long, complicated, difficult to verify, and fragile in the face of specification change.
- Automatic generation with mechanised verification is the future!

Future directions:

- Expand the session description language
- Finer-grain attacker-model
- Expand the modularity of the formal proof
- Towards matching existing protocols

Papers

- [CSF'07] [TGC'07] [CSF'09]
- Theoretical extension with concurrency [CONCUR'09]
- F* extends F7 [POPL'10] [POPL'12]

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